

Consistent Query Answering in Partially Consistent Databases

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A database is called uncertain if two or more tuples of the same relation are allowed to agree on their primary key. Intuitively, such tuples act as alternatives for each other. A repair (or possible world) of such uncertain database is obtained by selecting a maximal number of tuples without ever selecting two tuples of the same relation that agree on their primary key. For a Boolean query \mathbf{q} , the problem **CERTAINTY**(\mathbf{q}) takes as input an uncertain database \mathbf{db} and asks whether \mathbf{q} evaluates to true on every repair of \mathbf{db} . In recent years, the complexity of **CERTAINTY**(\mathbf{q}) has been studied under different restrictions on \mathbf{q} . These complexity studies have assumed no restrictions on the uncertain databases that are inputted to **CERTAINTY**(\mathbf{q}). In practice, however, it may be known that these input databases are partially consistent, in the sense that they satisfy some dependencies (e.g., functional dependencies).

In this talk, we introduce the problem **CERTAINTY**(\mathbf{q}) in the presence of a set Σ of dependencies. The problem **CERTAINTY**(\mathbf{q}, Σ) takes as input an uncertain database \mathbf{db} that satisfies Σ , and asks whether every repair of \mathbf{db} satisfies \mathbf{q} . We focus on the complexity of **CERTAINTY**(\mathbf{q}, Σ) when \mathbf{q} is an acyclic conjunctive query without self-join, and Σ is a set of functional dependencies and join dependencies, the later of a particular form. We provide an algorithm that, given \mathbf{q} and Σ , decides whether **CERTAINTY**(\mathbf{q}, Σ) is first-order expressible. Moreover, we show how to effectively construct a first-order definition of **CERTAINTY**(\mathbf{q}, Σ) if it exists.

References

- [1] Sergio Greco, Fabian Pijcke, Jef Wijsen, *Consistent Query Answering in Partially Consistent Databases*. Submitted 2013.